# Nutmeg: a MIP and CP Hybrid Solver Using Branch-and-Check 

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Decomposition at 70: In Honor of Jacques Desrosieres and His Legacy
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## Overview

- A generic method to find combinatorial Benders cuts.
- Based logic-based Benders decomposition.
- Generalizes classical Benders decomposition to discrete subproblems.
- No specific form of cuts due to its generality. User must derive cuts valid for their problem.
- Otherwise, naive combinatorial Benders cuts: $x_{1}, \ldots, x_{n}$ binary, $x_{1}+x_{2}+\ldots+x_{n} \leq n-1$
- Uses propagation (aka inference) and conflict analysis (aka conflict-driven clause learning) from constraint programming and Boolean satisfiability (SAT).
- Central idea: find a set of values to some master-problem variables that implicate infeasibility in the subproblem.


## Example <br> Vehicle Routing Problem with Time Windows

- Familiar problem to illustrate the method.
- Obviously not the best way to solve the VRPTW.
- Master problem has arc variables and network flow constraints.
- Travel time and time window constraints moved into the Benders subproblem.
- Ignore load in this example.

Implication Graph
$\mathrm{t}_{\mathrm{i}}$ : earliest time of starting service at vertex i $\mathrm{c}_{\mathrm{i}, \mathrm{j}}$ : cost/travel time from ito j

Network


Branch-and-Bound Tree

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Data
Decision (branching)
Propagation (inferred)


Network


Branch-and-Bound Tree

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## Implication Graph

$\mathrm{t}_{\mathrm{i}}$ : earliest time of starting service at vertex i $\mathrm{c}_{\mathrm{i}, \mathrm{j}}$ : cost/travel time from i to j

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Decision (branching)
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## $\mathrm{X}_{6,2}=1$

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## Implication Graph

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Branch-and-Bound Tree


## Implication Graph

$\mathrm{t}_{\mathrm{i}}$ : earliest time of starting service at vertex i $\mathrm{c}_{\mathrm{i}, \mathrm{j}}$ : cost/travel time from i to j
Data Decision (branching) Propagation (inferred)


Network


Branch-and-Bound Tree


## Caveat

- In the example: trace the propagations backwards in the implication graph until all literals concern variables in the master problem.
- In practice: trace the implication graph backwards until (1) all literals concern variables in the master problem and (2) there is only one literal at the current depth of the branch-and-bound tree. Get a "first unique implication point" (1UIP) nogood. Experimental evidence from SAT suggests 1UIP nogoods perform better.


## Nutmeg

- A proof-of-concept automatic decomposition solver that implements the generic form of logic-based Benders decomposition.
- Calls SCIP for the MIP master problem and branch-and-bound.
- Calls Geas for the CP Benders subproblem and conflict analysis.
- Actually, just a bunch of SCIP plug-ins that glue the master problem and subproblem together.


## Vehicle Routing Problem with Location Congestion

- Multiple customers have the same ( $x, y$ ) coordinates. Deliveries to the same warehouse.
- Warehouses have a limited number of machines for unloading the vehicles.
- Vehicles must share the machines to unload.
- VRPLC has joint routing and scheduling combinatorial structure (a type of sychronization).


## Experimental Results

- MIP is good at network flow (TSP, VRP).
- CP is good at packing (RCPSP).
- Hybrid appears to be good at vehicle routing with scheduling.



## Experimental Results

- Nutmeg tested on a range of problems.
- LBBD known to perform well on some problems (e.g., VRPLC).
- LBBD performance unknown on majority of problems. No one has tried due to labor of deriving problem-specific cuts.
- Results confirm LBBD successes without manually deriving cuts.
- Mostly does not work. Problem still needs appropriate structure.
- Discovered one new problem suitable for LBBD.


## Key Points

- Logic-based Benders decomposition previously required problem-specific cuts.
- Otherwise, naive combinatorial Benders cuts work but are very weak.
- Conflict analysis can find tighter combinatorial Benders cuts (fewer variables).
- Recently implemented in an automatic decomposition solver named Nutmeg.
- Nutmeg works only on problems with appropriate structure.
- In VRP, "robust" cuts over arcs naturally translate to cuts over paths in branch-and-price.
- VRPs with synchronization become easy/easier. Generate the routes independently and then prevent a subset of arcs that violate the synchronization constraints. But no guarantee it's fast.
- Super pre-alpha release at https://github.com/ed-lam/nutmeg.

